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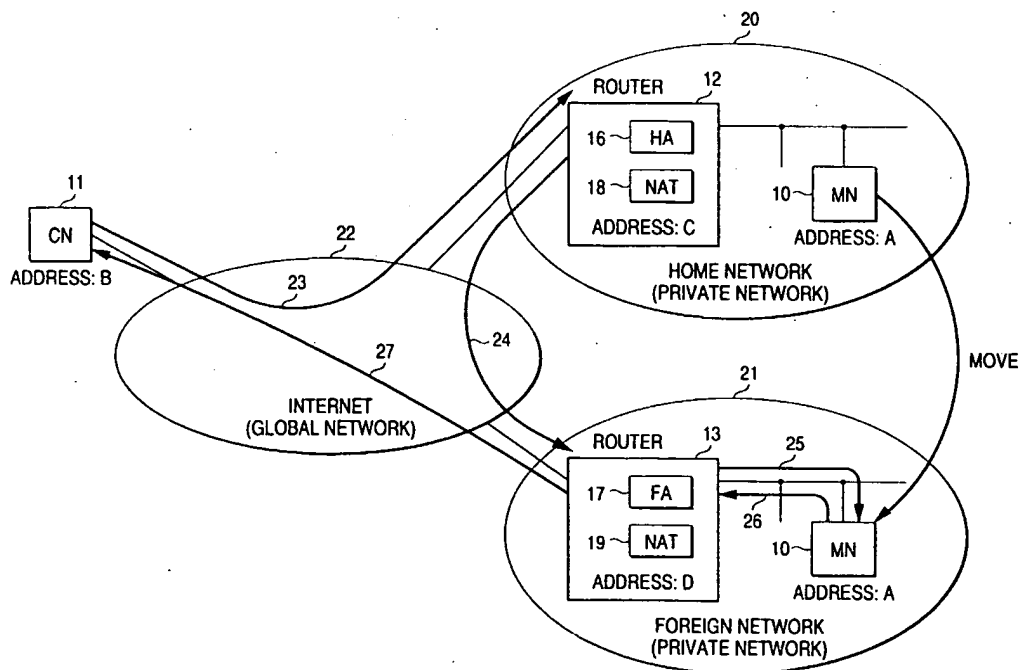
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(54) **Method for packet communication with mobile node**

(57) Even when a node (10) which can be moved among different networks (20, 21) and which sends a packet including a private address as a sender address sends the packet in either of the networks (20, 21), the sender address of the packet sent by the node (10) is

changed to a same global address. The method can prevent that a correspondent node (11) determines a packet from a same node as a packet from a different node due to a difference in their sender addresses, which hardly interrupts communication.

FIG. 6



Description

BACKGROUND OF THE INVENTION

1. Field of the Invention

[0001] The present invention relates to a method for packet communication where a node moves between different networks, and particularly to a method for packet communication where a node is moved to networks having a network address translation (NAT) function.

2. Description of the Related Art

[0002] In Internet Protocol (IP) communication, a wide area network is divided into a plurality of sub-networks (local area network) for management. Further, each node (node) is assigned an IP address, which includes a network portion indicating a number for identifying a sub-network and a host portion for indicating a reference number of a node in the sub-network.

[0003] In IP communication, routing is performed by including the IP address as a sender address and a destination address. For example, a router provided in each sub-network receives a packet existed on an IP network and checks whether a network portion of the destination address matches with a sub-network number managed by itself. If they match, the router incorporates the packet to the its managing sub-network. Each node selects and receives a packet whose destination address matches with its IP address among packets incorporated to the sub-network. It should be noted that the router might be called gateway in software terms.

[0004] Due to the above-described configuration of the IP network, a node to which an IP address is assigned within a particular sub-network cannot receive packets addressed to itself when the node is moved to another sub-network. It is because the router managing the sub-network to which the node is moved does not incorporate packets addressed to the node to its managing network.

[0005] Thus, in order to solve such a problem, a mobile IP (RFC2002, "IP Mobility Support", RFC: Request For Comment) has been proposed. The mobile IP is proposed by Internet Engineering Task Force (IETF). It is a technology for transferring an IP packet to a node without changing the IP address assigned to the node even when the node is moved across different sub-networks.

[0006] FIG. 1 is a construction diagram exemplifying a network concept of the mobile IP. FIG. 1 includes a mobile node (MN) 10, a home network 20 which is sub-network to which the mobile node 10 originally belongs, a home agent unit (HA) 16 for managing movements of the mobile node 10, a router 14 having the HA 16, a foreign network 21 which is a sub-network to which the mobile node 10 moves, a foreign agent unit (FA) 17 for managing the mobile node 10 which is moved into the foreign network 21, a router 15 having the FA 17, a cor-

respondent node (CN) 11 for communicating with the mobile node 10, and an IP network 22 such as Internet connected to the home network 20 and the foreign network 21.

[0007] A processing procedure of the mobile IP will be described by using as an example a case where the mobile node 10 is moved to the foreign network 21. An IP address is preset in the mobile node 10 in hardware or software manner. The IP address of the mobile node 10 includes an identification number of the home network 20 as a network portion and includes an identification number of the mobile node 10 in the home network 20 as a host portion.

[0008] Now, it is assumed that a packet addressed to the mobile node 10 is sent from the correspondent node 11 onto the IP network 22. The destination address of the packet, that is, the IP address of the mobile node 10 includes the identification number of the home network 20. Thus, the packet is incorporated to the router 14 managing the home network 20 via a path 23. The HA 16 having the router 14 recognizes in advance that the mobile node 10 has been moved to the foreign network 21 and transfers the packet to the router 15 managing the foreign network 21. The FA 17 having the router 15 incorporates the transferred packet to its managing sub-network, that is, the foreign network 21. The mobile node 10 receives the packet incorporated to the foreign network 21 and determines that the destination address matches with its address in order to receive the packet.

[0009] It should be noted that a tunneling technology is used for transferring packets from the HA 16 to the FA 17. The tunneling technology adds a header including a new destination address and a transferring address to a packet when transferring. The processing that a transferring node adds the header is called encapsulation processing while the processing that a destination node removes the header is called decapsulation processing. In this case where a packet is transferred from the HA 16 to the FA 17, an IP address D, as a transferring address, of the router 15 having the FA 17 and an IP address C, as a destination address, of the router 14 having the HA 16 are added as a header.

[0010] It will be described a case where the mobile node 10 exists in the foreign network 21 and sends a packet to the correspondent node 11. The mobile node 10 creates a packet including an IP address B of the correspondent node 11 and its IP address (that is, home address A) as the destination and sender addresses, respectively. The created packet is sent to the correspondent node 11 via the path 26, the router 15, and the path 27.

[0011] FIG. 2 is a sequence diagram for exemplifying a processing sequence of a mobile IP in a case where the mobile node 10 is moved from the home network 20 to the foreign network 21. In FIG. 2, 1000 to 1002 are processing where the mobile node 10 exists in the home network 20 while 1003 to 1015 are processing where the mobile node 10 exists in the foreign network 21.

[0012] First of all, the mobile node 10 in the home network 20 receives an agent advertisement from the router 14 so that it recognizes that a sub-network where it is positioned is managed by the HA 16 of the router 14 (processes 1000 to 1002). FIG. 3 is an explanatory diagram for exemplifying a format of the agent advertisement. As shown, the agent advertisement includes an IP address (Care of Address) of an agent node (that is, a node having the HA and FA) managing the sub-network. It should be noted that the agent advertisement is a packet in which a router discovery of an Internet control message protocol (ICMP) is extended and broadcasted in the sub-network.

[0013] Next, the mobile node 10 is moved from the home network 20 into the foreign network 21 (process 1003). In the foreign network 21, the mobile node 10 receives an agent advertisement from the router 15 and recognizes that it is moved into the sub-network managed by the FA 17 of the router 15 (process 1004). Then, the mobile node 10 requests the FA 17 of the router 15 to register its presence (process 1005). More specifically, a message in a format shown in FIG. 4 is sent as a registration request message to the router 15. As shown in FIG. 4, the registration request message includes a home address of the mobile node 10, an address of a home agent managing the mobile node 10 (IP address of the router 14, here), and an address of a foreign agent (Care of address, IP address of the router 15, here).

[0014] The FA 17 receives the registration request message shown in FIG. 4 from the mobile node 10 and registers (stores) its content. Then, the FA 17 transfers the registration request message to the HA 16 (process 1006). The HA 16 receives the registration request message and registers (stores) its registration content so that packet communication can be performed in the foreign network 21 as well. Further, the HA 16 responds to the FA 17 that registration has been completed. More specifically, the FA 17 sends a message in a format shown in FIG. 5 as a response message to the router 15. As shown in FIG. 5, the message includes a home address of the mobile node 10 and an IP address of the home agent (IP address of the router 14, here) FA 17 receives the response message and sends the response message to the mobile node 10. The mobile node 10 receives the response message so that it can recognize that communication is possible (process 1008).

[0015] Next, a method for communication between the mobile node 10 and the correspondent node 11 will be described with reference to FIGS. 1 and 2. The mobile node 10 creates packet data whose destination address is an IP address B of the correspondent node 11 and whose sender address is its IP address A. Then, the mobile node 10 sends the packet into the foreign network 21 (process 1009). The packet is sent to the correspondent node 11 on the IP network 22 through the router 15 (process 1010). When the correspondent node 11 responds to the received packet, that is, when

the correspondent node 11 responds to the mobile node 10, the correspondent node 11 sends to the IP address B a packet whose destination address is the address A which is a destination address of the received packet and whose sender address is its IP address B (process 1011).

[0016] The packet sent in this manner includes an identification number of the home network 20 in the network portion of its destination address A. Thus, it is incorporated to the router 14 once. The home agent 16 of the router 14 has already realized that the mobile node 10 was moved to the foreign network 21. The home agent 16 encapsulates the received packet and tunneling-transfers it to the router 15 (processes 1012 and 1013). More specifically, the home agent 16 adds to the packet a new packet whose destination address is an IP address D of the router 15 and whose sender address is an IP address C of the router 14 and sends out the packet to the IP network 22. The router 15 receives the tunnel-transferred and encapsulated packet and decapsulates the packet (process 1014). It should be noted that the destination and sender addresses of the packet resulted from the decapsulation are the same as those before encapsulation, and they are the IP address A of the mobile node 10 and the IP address B of the correspondent node 11, respectively. The decapsulated packet are sent from the router 15 into the foreign network 21 and received by the mobile node 10. In this way, the communication is possible between the mobile node 10 in the foreign network 21 and the correspondent node 11 on the IP network.

[0017] Recently, since IP address resources must be used efficiently, a private network has been established which uses technologies including private addresses, NAT, and IP masquerades in enterprises (NAT: Network Address Translation).

[0018] It should be noted that when an address used within a particular network is a global address, a private address is an address used only in a particular sub-network belonging to the global network. Further, NAT translates a destination address and a sender address of each packet from a private address to a global address, or from a global address to a private address when packet communication is performed between node using the private address and node using the global address. In NAT, the private address and the global address are mapped one-to-one.

[0019] The IP masquerade is resulted from the extension of the NAT. The IP masquerade is a technology for mapping a plurality of private addresses to one global address by regarding one including a port number of TCP (transmission control) and/or UDP (user datagram protocol) as an address. A network using a private address and a network using a global address are called a private network and a global network, respectively, below.

SUMMARY OF THE INVENTION

[0020] It is an object of the present invention to provide a method for communication which can be applied to a case where such a private network is included in networks where a node is moved.

[0021] In order to achieve the object, the present invention changes a sender address of a packet sent by a node to a same global address when the node, which can be moved among different networks and sends a packet including a private address as a sender address, sends the packet in either one of the networks.

[0022] According to the construction, a sender address of a packet received in a correspondent node does not change before and after the movement of the node. Therefore, it reduces interruptions of communication due to that the correspondent node determines a packet from a same node as a packet from a different node.

BRIEF DESCRIPTION OF THE DRAWINGS

[0023]

FIG. 1 is a configuration diagram for exemplifying a communication system for a mobile IP;

FIG. 2 is a diagram for exemplifying a communication sequence of the mobile IP;

FIG. 3 is a diagram for exemplifying a format for an agent advertisement message used in the mobile IP;

FIG. 4 is a diagram for exemplifying a format for a registration request message used in the mobile IP; FIG. 5 is a diagram for exemplifying a format for a registration response message used in the mobile IP;

FIG. 6 is a configuration diagram for exemplifying a communication system according to a first embodiment of the present invention;

FIG. 7 is a diagram for exemplifying a communication sequence if registration and data transfer according to the first embodiment of the present invention;

FIG. 8 is a diagram for exemplifying address translation information according to the first embodiment of the present invention;

FIG. 9 is a diagram for exemplifying a format for a message notifying address translation information according to a first embodiment of the present invention;

FIG. 10 is a diagram for exemplifying a communication sequence according to a second embodiment of the present invention;

FIG. 11 is a diagram for exemplifying a format for a message requesting address translation information according to the second embodiment of the present invention;

FIG. 12 is a diagram for exemplifying a communi-

cation sequence according to a third embodiment of the present invention;

FIG. 13 is a diagram for exemplifying a communication sequence according to a fourth embodiment of the present invention;

FIG. 14 is a construction diagram for exemplifying a communication system according to a fifth embodiment of the present invention;

FIG. 15 is a diagram for exemplifying a communication sequence according to a fifth embodiment of the present invention;

FIG. 16 is a diagram for exemplifying a communication sequence according to a sixth embodiment of the present invention;

FIG. 17 is a diagram for exemplifying a communication sequence according to a seventh embodiment of the present invention;

FIG. 18 is a construction diagram for exemplifying a communication system according to a eighth embodiment of the present invention;

FIG. 19 is a diagram for exemplifying a communication sequence according to the eighth embodiment of the present invention; and

FIG. 20 is a diagram for exemplifying an address translation list according to the eighth embodiment of the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

[0024] When a mobile node is moved from a private network to a private network, from a private network to a global network, or from a global network to a private network, communication may be interrupted due to some setting in a correspondent node.

[0025] For example, in a case where a mobile node is moved from a private network to a private network, that is, where both home network and foreign network are private networks, and when the mobile node exists in the home network, address translation is performed by a NAT in the home network. When the mobile node is moved to the foreign network, the address translation is performed by another NAT in the foreign network. So, an IP address resulted from the translation by the NAT in the home network and an IP address resulted from the translation by the NAT in the foreign network are not translated into the same address.

[0026] Therefore, when the IP address resulted from the translation by the NAT in the private network the mobile node is moved from and the IP address resulted from the translation by the NAT in the private network the mobile node is moved to are different, the correspondent node determines those packets as packets from different mobile nodes, which causes interruption of communication.

[0027] Thus, according to the present invention, the communication is performed in a following manner.

First Embodiment

[0028] A method for packet communication according to the present invention will be described below with reference to figures. FIG. 6 is an explanatory diagram for exemplifying a network construction and its packet transfer path according to a first embodiment of the present invention. In FIG. 6, identical reference numerals will be given to identical or corresponding parts, and their description will be omitted here. In FIG. 6, both home network 20 and foreign network 21 are private networks using private addresses. A router 12 has functions of HA 16 and a NAT 18 and connects an IP network 22 and a home network 20. A router 13 has functions of the FA 17 and the NAT 19 and connects the IP network 22 and the foreign network 21. Assigned to a home address of the mobile node 10, the correspondent node 11, the router 12, and the router 13 are a private address A, a global address B, a global address C, and a global address D. While a case where the NAT 18 assigns the global address C of the router 12 itself to the mobile node 10, another global address may be assigned to the mobile node 10.

[0029] As described above, the NAT maps a particular global address to a private address of a node. Such a mapping relationship will be referred to address translation information below. FIG. 8 is an explanatory diagram for exemplifying address translation information for the NAT 18. FIG. 8 includes a private address 28 which is assigned to the mobile node 10 as a home address, and a global address 29 which is assigned as a translated address of the home address. Here, a global address C is mapped to a private address A. That is, the NAT 18 translates a sender address of a packet sent from the mobile node 10 from the private address A to the global address C. Further, the NAT 18 translates a destination address of a packet sent from the correspondent node 11 from the global address C to the private address A.

[0030] The first embodiment is characterized in that the NAT 19 of the foreign network 21 has identical address translation information to that of the NAT 18 in the home network 20 with respect to the mobile node 10. More specifically, the NAT 18 and the NAT 19 communicates each address translation information with each other periodically so that the identical global address are mapped for the mobile node 10. Thus, like the NAT 18, the NAT 19 translates a sender address of a packet sent from the mobile node 10 from the private address A to the global address C in order to send it to the correspondent node 11. Therefore, sender addresses of the packets sent from the home network 20 and the foreign network 21, respectively, can be matched, which can improve interruption of communication due to changes in sender addresses during communication.

[0031] A packet in a format shown in FIG. 9 is used for notifying address translation information. FIG. 9 shows message type 30, which indicates that it is a mes-

sage for notifying address translation information here. Further, FIG. 9 shows an unused reserved area, a number of address translation information to be notified, a home address 33 (for example, private address A) of the mobile node 10, and a global address 34 (for example, global address C) mapped to the home address of the mobile node 10. the home address 33 and the global address 34 make a pair, and a number of this pair is used as the number of the address translation information 32.

[0032] FIG. 7 is an explanatory diagram for exemplifying a sequence in a case where the mobile node 10 is moved from the home network 20 to the foreign network 21 during communication with the correspondent node 11 according to the network construction shown in FIG. 6. FIG. 7 includes processes 100 to 105 in a case where the mobile node 10 exists in the home network 20, and processes 107 to 123 in a case where the mobile node 10 exists in the foreign network 21. In the process 100 in FIG. 7, the mobile node 10 in the home network 20 sends packet data to the router 12 so that it can send data to the correspondent node 11 (process 100). A destination address of the packet is the global address B of the correspondent node 11, while its sender address is the private address A of the mobile node 10. It should be noted that packets having the destination address B and the sender address A as referred to Data (B, A) in FIG. 7. The other packets are similarly referred to Data (destination address, sender address).

[0033] The router 12 receives the packet from the mobile node 10. Then, the router 12 uses the NAT 18 to translate the sender address of the packet from the private address A to the global address C of the router 12 (process 101). The packet whose address has been translated is sent onto the IP network 22 and received by the correspondent node 11 (process 102). The correspondent node 11 sends the packet to the sender address of the packet received in the above-described process 102 in order to return the packet to the mobile node 10 (process 103). That is, the correspondent node 11 sends onto the IP network 22 the packet whose destination address is the global address C of the router 12 and whose sender address is the global address B of the correspondent node 11.

[0034] The router 12 receives the packet returned from the correspondent node 11. Then, the router 12 uses the NAT 18 to translate the destination address of the packet from the global address C of the router 12 to the private address A of the mobile node 10 (process 104). The packet having translated address is sent onto the home network 20 and received by the mobile node 10 (process 105). This way allows data exchanges between the mobile node 10 in the home network 20 and the correspondent node 11 on the IP network.

[0035] Next, it will be described a case where the mobile node 10 is moved from the home network 20 to the foreign network 21 during communication as described above. In the process 106 in FIG. 7, the mobile node 10

is moved from the home network 20 to the foreign network 21 (process 106). The mobile node 10 receives an agent advertisement in the foreign network 21 so that the mobile node 10 can detect that it is moved to the foreign network 21 (process 107). As described with reference to FIG. 3, the agent advertisement includes an IP address of the FA 17 (IP address D of the router 13, here).

[0036] The mobile node 10 having detected the movement requests registration to the FA 17 by using a format shown in FIG. 4 as described above. The FA 17 receives the registration request and notices that the mobile node 10 has been moved into the foreign network 21. Further, the FA 17 transfers the registration request to the HA 16 (process 109). The HA 16 receives the registration request from the FA 17 and notices that the mobile node 10 has been moved to the foreign network 21. Then, the HA 16 returns a registration response in a format shown in FIG. 5 described above to the FA 17 (process 110). Further, the registration response is returned from the FA 17 to the mobile node 10 so that the registration will be completed (process 111).

[0037] After the registration process as described above, the NAT 18 and the NAT 19 notify address translation information regarding the mobile node 10 to each other periodically. It will be described in detail. The NAT 18 has a timer, which times out periodically (process 112). Corresponding to the time-out, address translation information is notified from the NAT 18 to the NAT 19 in a format shown in FIG. 9 described above (process 113). Thus, the NAT 19 notices that the home address A of a packet of the mobile node 10 is translated to the global address C in the home network 20. Similarly, when a timer provided in the NAT 19 times out (process 114), address translation information is notified from the NAT 19 to the NAT 18 (process 115).

[0038] Next, a case will be described where data is sent from the mobile node 10 in the foreign network 21 to the correspondent node 11. The mobile node 10 in the foreign network sends a packet to the router 13 in order to send data to the correspondent node 11 (process 116). A destination address and sender address of the to-be-sent packet are the global address B of the correspondent node 11 and the private address A of the mobile node 10, respectively. The NAT 19 of the router 13 having received the packet from the mobile node 10 uses address translation information received in process 113 to translate a sender address of the packet from the private address A to the global address C. The destination address and the sender address of the packet whose address is translated in this manner are both global addresses. Therefore, the packet can be sent onto the IP network and transmitted to the correspondent node 11 (process 118).

[0039] The correspondent node 11 having received the packet send the packet whose destination address is the global address C which is the sender address of the earlier received packet in order to return the packet

to the mobile node 10 (process 119). The router having received the packet from the correspondent node 11 translates the destination address of the packet from the global address C to the private address A in the NAT 18 (process 120). Further, HA 16 adds a new header to the packet so that the destination address and the sender address become the global address D of the router 13 and the global address C of the router 12, respectively (that is, IP-encapsulated). The encapsulated packet is translated, that is, tunneled to the FA 17 of the router 13 (process 121). The FA 17 of the router 13 having received the tunneled packet decapsulates the packet (process 122) and sends it to the mobile node 10 (process 123). This allows communication between the mobile node 10 in the foreign network 21 and the correspondent node 11 on the IP network.

[0040] As described above, according to the first embodiment, even when a mobile node to which a private address is assigned is moved between different sub-networks, interruption of communication due to changes in a sender address during communication can be improved since mapping between a private address and a global address with respect to the mobile node in a foreign network is the same as mapping in a home network.

Second Embodiment

[0041] In the first embodiment, as shown in the processes 112 to 115 in FIG. 7, address translation information is sent by using time-out of a timer as a trigger. In a second embodiment, a registration request to be sent from the mobile node 10 to the router 13 is used as a trigger in order to send a request message for address translation information from the NAT 19 to the NAT 18. FIG. 10 is an explanatory diagram for exemplifying a sequence where, in a network construction shown in FIG. 6, a registration request from the mobile node 10 is used as a trigger so that the NAT 19 of the router 13 obtains address translation information from the NAT 18 of the router 12. FIG. 10 includes each of processes 200 to 207.

[0042] In process 200 in FIG. 10, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 200). The mobile node 10 receives an agent advertisement so that it detects that it has been moved to the foreign network 21 (process 201). The mobile node 10 having detected the movement uses a format shown in FIG. 4 in order to request a registration to the FA 17 (process 202). The receipt of the translation request is used as a trigger, and the NAT 19 of the router 13 requests address translation information regarding the mobile node 10 to the NAT 18 of the router 12 (process 203). FIG. 11 is an explanatory diagram for exemplifying a format of a message sent to the NAT 18 as the address translation information request. 40 indicates a message type, which includes a code indicating that it is a message for requesting address translation information here. 41 indicates that a reserved area is un-

used, and 42 indicates a number of address translation information to be requested. 43 indicates a home address of the mobile node 10.

[0043] The NAT 18 having received the address translation information request in the process 203 returns address translation information regarding the mobile node 10 to the NAT 19 (process 204). That is, a global address mapped to the home address of the mobile node 10 (a private address here) is sent to the NAT 19 by using a format shown in FIG. 9. Next, a registration is requested from the FA 17 to the HA 16 (process 205) and a response indicating that a registration has been done, that is a registration response is returned from the HA 16 to the FA 17 (process 206). Further, the registration response is returned from the FA 17 to the mobile node 10 (process 207).

[0044] In the manner as described above, the mobile node 10 is registered to the HA 16 and the FA 17, and address translation information of the mobile node 10 is set in the NAT 19. Thus, a sender address of a packet sent from the mobile node 10 is translated to a global address, which is identical to that of the NAT 18 of the router 12, in the NAT 19 of the router 13.

Third Embodiment

[0045] In the second embodiment, a registration request sent from the mobile node 10 to the router 13 is used as a trigger in order to request address translation information. In a third embodiment, a registration response sent from the NAT 18 to the NAT 19 is used as a trigger in order to request address translation information from the NAT 19 to the NAT 18. FIG. 12 is an explanatory diagram for exemplifying a sequence where, in a network construction shown in FIG. 6, a registration response from the HA 16 is used as a trigger so that the NAT 19 of the router 13 obtains address translation information from the NAT 18 of the router 12. FIG. 12 includes each of processes 300 to 307. In process 300 in FIG. 12, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 300). The mobile node 10 receives an agent advertisement so that it detects that it has been moved to the foreign network 21 (process 301). The mobile node 10 having detected the movement uses a format shown in FIG. 4 in order to request a registration to the FA 17 (process 302). The FA 17 having received the registration request requests a translation further to the HA 16 after the registration (process 303). The HA 16 having received the registration request from the FA 17 returns a response indicating that it has registered, that is, a registration response to the FA 17 after the registration (process 304).

[0046] When the router 13 receives the registration response in the process 304 from the HA 16, it requests address translation information to the NAT 18 of the router 12 by using the receipt of the registration response as a trigger (process 305). More specifically, a message requesting address translation information re-

garding the mobile node 10 is sent by using a format shown in FIG. 11. The NAT 18 having received the request message returns the address translation information by using a format shown in FIG. 9 to the NAT 19 (process 306). The NAT 19 of the router 13 receives the address translation information and then returns a translation response to the mobile node 10 (process 307).

[0047] As described above, address translation information is requested after an HA (home agent) returns a translation response, which eliminates a need for address translation information requests when a registration is failed. In this case, an unnecessary message does not have to be sent to a network.

Fourth Embodiment

[0048] In a fourth embodiment, the address translation information request shown in FIG. 11 is added into a message of the registration request and the registration response in the processes 109 and 110 shown in FIG. 7. FIG. 13 is an explanatory diagram for exemplifying a sequence where, in a network construction shown in FIG. 6, an address translation information request is added into a translation request message and address translation information is added into a translation response message so that the NAT 10 of the router 13 obtains address translation information from the NAT 19 of the router 12.

[0049] In a process 400 in FIG. 13, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 400). The mobile node 10 receives an agent advertisement so that it detects that it has moved to the foreign network 21 (process 401). The mobile node 10 having detected the movement requests a registration to the FA 17 by using the format shown in FIG. 4 (process 402). The FA 17 having received the translation request requests an extended registration to the HA 16 after translation. More specifically, the FA 17 sends an extended registration request message, which is made by adding a number of address translation information requests 42 and a home address 43 shown in FIG. 11 to the format shown in FIG. 4 (process 403).

[0050] The HA 16 having received the extended registration request in the process 403 returns an extended registration response to the HA 17. More specifically, the HA 16 returns an extended registration response message, which is made by adding a number of address translation information 32, a home address 33, and a translation address 34 shown in FIG. 9 to the format shown in FIG. 5 (process 404). The NAT 19 receives the extended registration response so that it can obtain address translation information regarding the mobile node 10. Then, a registration response is returned from the FA 17 to the mobile node 10 so that the registration is completed (process 405).

[0051] As described above, an address translation information request message is added to a registration request message for transmission. Thus, it can reduce a

number of message sent and time required from the registration request to a completion of obtaining address translation information.

Fifth Embodiment

[0052] In the first embodiment, both home network 20 and foreign network 21 have NATs. In a fifth embodiment, the foreign network 21 does not have a NAT. It will be described a case where the FA 17 performs address translation with respect the mobile node 10. FIG. 14 is an explanatory diagram for exemplifying a network construction and its packet transfer path according to the fifth embodiment. In FIG. 14, identical reference numerals will be given to identical or corresponding parts in FIG. 6, and their description will be omitted here.

[0053] In FIG. 14, the foreign network 21 is a global network using a global address. The router 15 connects an IP network 22 and a foreign network 21. Further, the home network 20 is a private network using a private address. The router 12 connects the IP network 22 and the home network 20. The router 15 has functions of the FA 17 but not NAT functions. The router 12 has functions of the HA 16 and the NAT 18. A private address A, a global address B, a global address C, and a global address D are assigned to a home address of the mobile node 10, the correspondent node 11, the router 12, and the router 15, respectively.

[0054] FIG. 15 is an explanatory diagram for exemplifying a sequence where, in a network construction shown in FIG. 14, the mobile node 10 is moved from the home network 20 to the foreign network 21 during communication with the correspondent node 11. In FIG. 14, 500 to 522 indicate processes, respectively. In the process 500 in FIG. 15, the mobile node 10 in the home network 20 sends a packet to the router 12 in order to send data to the correspondent node 11 (process 500). A destination address of the packet is a global address B while its sender address is a private address A.

[0055] The router 12 having received the packet from the mobile node 10 uses the NAT 18 to translate the sender address of the received packet from the private address to the global address C of the router 13 (process 501). Then, the packet whose address is translated is sent to the correspondent node 11 (process 502). The correspondent node 11 receives the packet and returns the packet whose destination address is the address C of the router 12 and whose sender address is the address B of the correspondent node 11 (process 503). The router 12 having received the returned packet from the correspondent node 11 uses the NAT 18 to translate the destination address of the returned packet from the address C to the address A of the mobile node 10 (process 504). The packet whose address has been translated is sent to the mobile node 10 (process 505). In this way, communication is performed between the mobile node 10 in the home network 20 and the correspondent node 11 on the IP network.

[0056] Further, during communication as above, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 506). The mobile node 10 receives an agent advertisement in the foreign network 21 so that it detects that it has moved to the foreign network 21 (process 507). The mobile node 10 having detected the movement request a registration to the FA 17 by using the format shown in FIG. 4 (process 508). Further, a registration request is performed from the FA 17 to the HA 16 (process 509), and a registration response is returned from the HA 16 to the FA 17 (process 510). As shown in FIG. 5, the registration response includes a home address of the mobile node 10.

[0057] Here, the home address of the mobile node 10 is a private address. Thus, The FA 17 having received the registration response in the process 510 determines that address translation information regarding the mobile node 10 is required (process 511). Then, the FA 17 sends a message requesting for address translation information regarding the mobile node 10 to the NAT 18 by using the format shown in FIG. 11 (process 512). The NAT 18 returns address translation information of the mobile node 10 in the format shown in FIG. 9 to the FA 17 as a response to the request message in the process 512 (process 513). The FA 17 receives the response and then sends a registration response to the mobile node 10 (process 514). Thus, a registration is completed, and address translation information of the mobile node 10 is set in the FA 17.

[0058] Next, a case will be described where data is sent from the mobile node 10 in the foreign network 21 to the correspondent node 11. The mobile node 11 in the foreign network 21 sends a packet to the router 15 in order to send data to the correspondent node 11 (process 515). A destination address of the packet is a global address B while its sender address is a private address A. The router 15 having received the packet from the mobile node 10 translates through the FA 17 the sender address of the received packet from the private address A to the global address C based on address translation information obtained in the above-described process 513 (process 516). The packet whose address has been translated is sent to the correspondent node 11 (process 517).

[0059] The correspondent node 11 receives the packet and then sends as a returned packet to the mobile node 10 the packet whose destination address is an address C which is a sender address of the received packet (process 518). The NAT 18 of the router 12 translate the destination address of the returned packet from the correspondent node 11 from the address C of the router 12 to the address A of the mobile node 10. Further, through the HA 16, the NAT 18 sets the destination address and the sender address to the address D of the router 15 and the address C of the router 12 for IP-encapsulation (process 519). The IP-encapsulated packet is tunneled to the FA 17 of the router 15 (process 520). The FA 17 of the router 15 decapsulates the tunneled

packet (process 521). The packet resulted from the decapsulation is sent to the mobile node 10 (process 522). In this way, communication becomes possible between the mobile node 10 in the foreign network 21 and the correspondent node 11 on the IP network.

[0060] As described above, when an NAT does not exist in a foreign network, an FA (foreign agent) performs address translation for a packet received/sent by a mobile node. Thus, the mobile node can continue communication even in the foreign network. That is, even when a home network is a private network and a foreign network to which it is moved is a global network, communication can be continued.

Sixth Embodiment

[0061] In the fifth embodiment, the FA 17 performs address translation with respect to the mobile node 10. In a sixth embodiment, the mobile node 10 itself performs address translation. Further, the mobile node 10 uses a registration response sent from the FA 17 as a trigger to obtain address translation information. FIG. 16 shows in the network construction shown in FIG. 14 a sequence where the mobile node 10 itself uses a registration response as a trigger to obtain address translation information from the NAT 18 of the router 12.

[0062] In a process 600 in FIG. 16, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 600). The mobile node 10 receives an agent advertisement so that it detects that it has moved to the foreign network 21 (process 601). The mobile node 10 having detected its movement uses the format shown in FIG. 3 to request a registration to FA 17 (process 602). After the registration, the FA 17 further requests a registration to the HA 16 (process 603). After the registration, the HA 16 returns a registration response to the FA 17 (process 604). The FA 17 receives the registration response and sends the registration response further to the mobile node 10 (process 605). In this way, the registration is completed.

[0063] The mobile node 10 having received the registration response from the FA 17 uses it as a trigger in order to send a message requesting its address translation information to the NAT 18 by using the format shown in FIG. 11 (process 606). The NAT 18 sends address translation information of the mobile node 10 to the mobile node 10 by using the format shown in FIG. 9 as a response to the request message (process 607). The mobile node 10 receives the response and notices its address translation information.

[0064] Next, it will be described a case where data is sent from the mobile node 10 in the foreign network 21 to the correspondent node 11. In process 608, the mobile node 10 creates packet data to be sent to the correspondent node 11. A sender address of the created packet here is not the private address A of the mobile node 10 but is set to the global address C which is mapped to the private address A. The created packet is

sent, and then it is received by the correspondent node 11 (process 609). The correspondent node 11 sends a packet by using the address C which is a sender address of the received packet for its destination address as a response to the received packet (process 610). The packet whose destination address is the address C is received by the router 12. The router 12 uses the NAT 18 in order to translate the destination address of the received packet from the global address C to the private address A (process 611). Further, the HA 16 of the router 12 recognizes by a registration request in the process 603 that the mobile node 10 is under a management network of the FA 17. The HA 16 IP-encapsulates the address translated packet (process 611) and tunneling-transfers it to the FA 17 of the router 15 (process 612). The FA 17 of the router 15 decapsulates the tunneling-transferred packet in order to obtain the original packet whose destination address is the private address A (process 613). The original packet obtained by the decapsulation is sent to the mobile node 10 (process 614). **[0065]** In this way, communication becomes possible between the mobile node 10 in the foreign network 21 and the correspondent node 11 on the IP network. Particularly, when the mobile node itself to which a private address is assigned obtains address translation information, communication is possible in a foreign network which is a global network without addition of functions to an FA (foreign network) as much as possible thereby.

Seventh Embodiment

[0066] In the sixth embodiment, the mobile node 10 uses a registration response as a trigger in order to obtain address translation information. In a seventh embodiment, a registration request sent from the mobile node 10 to the FA 17, a registration request sent from the FA 17 to the HA 16, a registration response sent from the HA 16 to the FA 17, and a registration response sent from the FA 17 to the mobile node 10 are extended respectively so that the mobile node 10 can obtain address translation information.

[0067] FIG. 17 shows a sequence where, in the network construction shown in FIG. 14, an address translation information request is added to a registration request message, and the address translation information is added to a registration response message, so that the mobile node 10 obtains address translation information from the NAT 18 of the router 12. In a process 700 in FIG. 17, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 700). The mobile node 10 receives an agent advertisement so that it can detect that it has moved to the foreign network 21 (process 701).

[0068] The mobile node 10 having detected the movement in the process 701 performs an extended registration request in order to request its address translation information to the FA 17 (process 702). More particularly, an extended registration request message which is

made by adding the number of address translation information requests 42 and the home address 43 shown in FIG. 11 to the format shown in FIG. 4 is sent to the FA 17. The FA 17 receives the extended registration request message in order to perform a predetermined registration. Then, the message is sent to the router 12 (process 703). The router 12 receives the message and performs a predetermined registration in the HA 16. Also, the extended registration response is sent to the FA 17 by including address translation information regarding the mobile node 10 in the message (process 704). More specifically, an extended registration response message by adding the address translation information 32, the home address 33, and the translation address 34 shown in FIG. 9 to the format shown in FIG. 5 is sent to the FA 17. the extended registration response message is further sent from the FA 17 to the mobile node 10 (process 705). Thus, the registration is completed, and address translation information regarding the mobile node 10 is set in the mobile node 10.

[0069] As described above, an address translation information request message is added to a registration request message for transmission, which can reduced a number and time for sending a message required from the registration request to obtaining the address translation information.

Eighth Embodiment

[0070] In the first embodiment, both home network 20 and foreign network 21 have NATs. In a eighth embodiment, it will be described a case where the home network 20 does not have a NAT while the foreign network 21 has the NAT. A global address is assigned to the mobile node 10 here. FIG. 18 is an explanatory diagram for exemplifying a network construction and its packet transfer path according to the eighth embodiment. In FIG. 18, identical reference numerals will be given to identical or corresponding parts in FIG. 6, and their description will be omitted here.

[0071] In FIG. 18, the home network 20 is a global network using a global address. The router 14 does not have a NAT function. More specifically, the router 14 has functions of the HA 16 and connects the IP network 22 and the home network 20. The foreign network 21 is a private network using a private address. The router 15 has functions of the FA 17 and NAT 19 and connects the IP network 22 and the foreign network 21. A global address A, a global address B, a global address C, and a global address D are assigned to a home address of the mobile node 10, the correspondent node 11, the router 14, and the router 13, respectively.

[0072] FIG. 19 is an explanatory diagram for exemplifying a sequence where, in a network construction shown in FIG. 18, the mobile node 10 is moved from the home network 20 to the foreign network 21 during communication with the correspondent node 11. In FIG. 18, 800 to 818 indicate processes, respectively. In the proc-

ess 800 in FIG. 18, the mobile node 10 in the home network 20 sends a packet to the router 14 in order to send data to the correspondent node 11. A destination address of the packet is a global address B while its sender address is a private address A. The router 14 receives the packet and send it to the correspondent node 11 (process 801). The correspondent node 11 sends the packet whose destination address is a global address A which is a sender address of the received packet as a response to the packet from the router 14 (process 802). The packet is received by the router 14 (process 802) and then is sent from the router 14 to the correspondent node 11 (process 803). In this way, communication is performed between the mobile node 10 in the home network 20 and the correspondent node 11 on the IP network.

[0073] During communication as above, the mobile node 10 is moved from the home network 20 to the foreign network 21 (process 804). The mobile node 10 receives an agent advertisement so that it detects that it has moved to the foreign network 21 (process 805). The mobile node 10 having detected the movement requests a registration to the FA 17 of the router 13 by using the format shown in FIG. 4. (process 806). FA 17 receives the registration request and further performs a similar registration request to the HA 16 (process 807). HA 16 returns a registration response to the registration request by using the format shown in FIG. 5 (process 808). According to the home address of the mobile node 10 included in the response, the FA 17 can recognize that a global address is assigned to the mobile node 10. That is, it recognizes that address translation is not performed by a NAT in the home network 20 and address translation is not necessary for the packet from the mobile node 10.

[0074] The FA 17 of the router 13 which has recognized that address translation is not necessary for the home address A of the mobile node 10 registers it in an address non-translated list (process 809). The address non-translated list is a list for registering IP addresses on which the NAT 19 of the router 13 perform address translation. FIG. 20 is an explanatory diagram for exemplifying the address non-translated list. Here, an address A is registered, and the NAT 19 does not perform NAT translation on the address A registered in the list. After that, a registration response is returned from the FA 17 to the mobile node 10 in order to complete the registration (process 810).

[0075] Next, it will be described a case where data is sent from the mobile node 10 in the foreign network 21 to the correspondent node 11. In a process 811, the mobile node 10 in the foreign network 21 sends a packet to the router 13 in order to send data to the correspondent node 11 (process 811). A destination address of the packet is a global address B while its sender address is a global address A. The NAT 18 of the router 13 having received the packet from the mobile node 10 does not perform address translation on the global address A

based on the address non-translated list shown in FIG. 20 (process 812). The non-translated packet is sent on to the IP network 22 and received by the correspondent node 11 (process 813).

[0076] The correspondent node 11 sets the address A which is a sender address of the earlier received packet as its destination address and transmits the packet in order to transmit the packet to the mobile node 10 (814). The HA 16 of the router 14 having received the packet from the correspondent node 11 sets address D of the router 13 and address C of the router 14 in the destination address and the sender address, respectively, and performs IP-encapsulation thereon (process 815). The IP encapsulated packet is tunneled to the FA 17 of the router 13 (process 816). The router 13 of the FA 17 decapsulates the tunneled packet while does not perform address translation through the NAT 19 based on the address non-translated list (process 817). Then, the decapsulated packet is sent to the mobile node 10 (process 818). In this way, communication becomes possible between the mobile node 10 in the foreign network 21 and the correspondent node 11 on the IP network.

[0077] The global addresses which are assigned to the router 12, HA 16, and NAT 18 in FIG. 6 has been described above as a single address C. However, different global addresses may be assigned, respectively. Similarly, different global addresses may be assigned to the router 13, the FA 17, and the NAT 19, respectively.

[0078] A method such as for example the above is stored on computer readable medium such as HARD DISK, FLOPPY DISK, CDROM, DVDROM, MEMORY.

Claims

1. A method for packet communication, wherein a sender address of a packet including a private address as a sender address sent by a node (10) which can be moved among different networks (20, 21) is changed to a same global address even when the node (10) sends the packet in either of the networks (20, 21).

2. A method for packet communication, wherein a node (10) including a private address as a sender address is moved from a first network (20) having the node (10) and address changing means for receiving the packet from the node (10) and then changing the sender address of the packet to a global address for output into a second network (21) which is different from the first network (20), comprising the following steps:

- receiving the packet from the node (10) which is moved into the second network (21);
- changing the sender address of the received packet from the private address to a same global address as the global address; and

- sending the packet whose sender address has been changed to the outside of the second network (21).

3. The method according to claim 2, comprising the following step:

the node (10) moved into the second network (21) and the address changing means in the first network (20) notify to each other correspondence between the private address and the global address periodically after the node (10) has been registered in a home agent (16) for managing the first network (20) and a foreign agent (17) for managing the second network (21).

4. The method according to claim 2 or 3, comprising the following steps:

- detecting that a registration request is sent from the node (10) moved into the second network (21) to a foreign agent (17) for managing the second network (21); and
- after the detection, requesting the address changing means in the first network (20) to change the global address mapped to the private address.

5. The method according to any of claims 2 to 4, comprising the following step:

- detecting that a response indicating that the node (10) is registered is sent from a home agent (16) for managing the first network (20) to a foreign agent (17) for managing the second network (21); and after the detection, requesting the address changing means in the first network (20) to change the global address mapped to the private address.

6. The method according to any of claims 2 to 5, further comprising the following step:

- adding a code for requesting the global address mapped to the private address to a registration request sent from a foreign agent (17) for managing the second network (21) to a home agent (16) for managing the first network (20).

7. A method for packet communication, comprising the following steps:

- outputting a packet including a private address as a sender address to changing means for changing a sender address of the packet from the private address to a global address; and
- sending a packet whose sending address is a

global address which is the same as the global address in a second network (21) which is different from a first network (20).

8. A method for packet communication, 5
comprising the following steps:
 - receiving a packet including a private address 10
and a sender address from a first node (10) in
a first network (20);
 - changing a sender address of the received 10
packet from the private address to a first global
address;
 - sending the packet whose sender address has 15
been changed to the outside of the first network
(20);
 - receiving a packet including a second global 20
address, which is different from the first global
address as a sender address, from a second
node (11) which has been moved into the first
network (20) from the second network (21)
which is different from the first network (20);
and
 - sending the packet to the outside of the first net- 25
work (20) without changing a sender address
of the packet received from the second node
(11) from the second global address to the first
global address.
9. A computer program stored on a computer readable 30
medium, wherein a node (10) including a private ad-
dress as a sender address is moved from a first net-
work (20) having the node (10) and address chang-
ing means for receiving the packet from the node 35
(10) and then changes the sender address of the
packet to a global address for output into a second
network (21) which is different from the first network
(20), comprising:
 - a code for receiving the packet from the node 40
(10) which is moved into the second network
(21);
 - a code for changing the sender address of the 45
received packet from the private address to a
same global address as the global address;
and
 - a code for sending the packet whose sender 50
address has been changed to the outside of the
second network (21).

50

55

FIG. 1

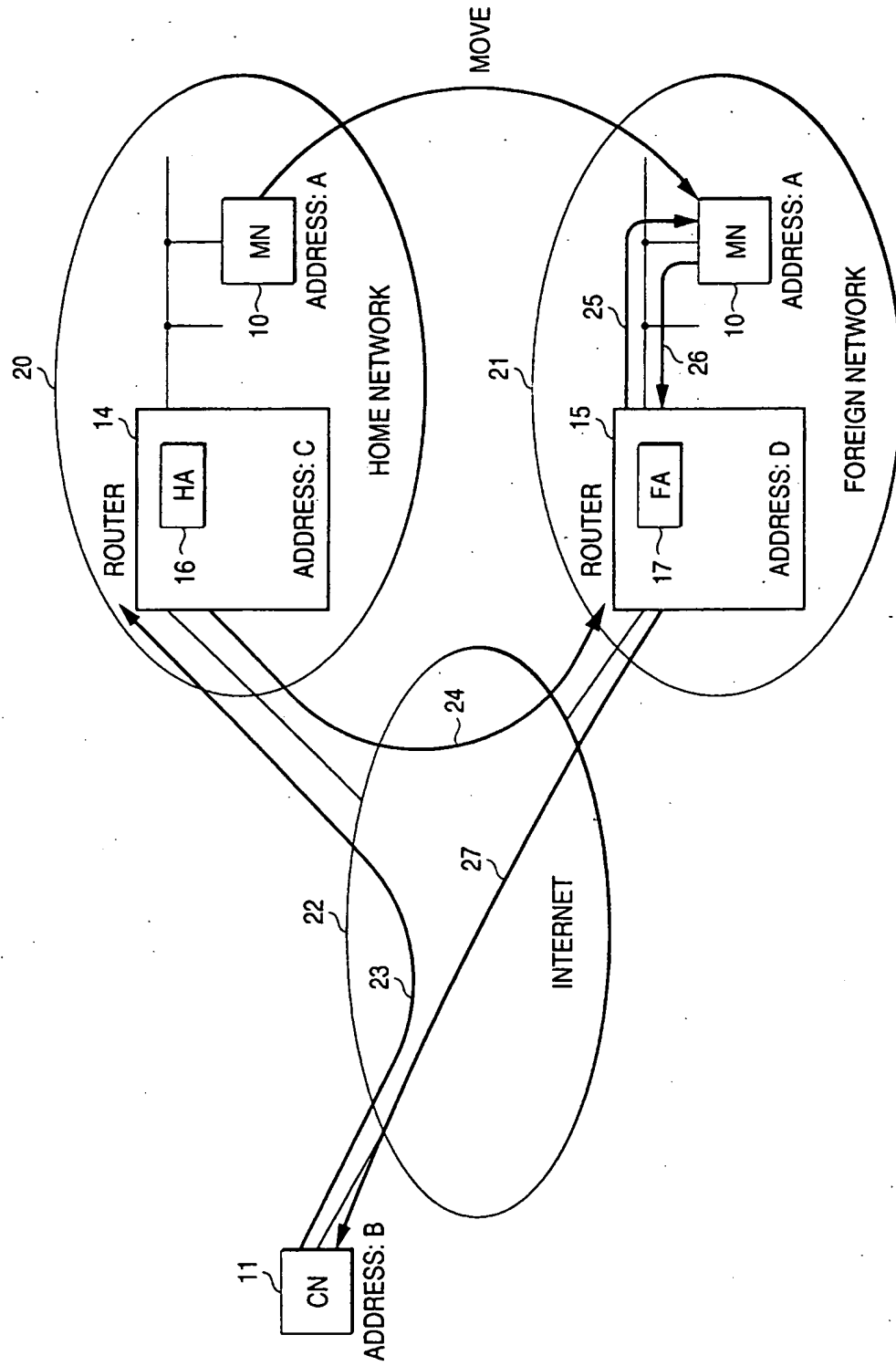


FIG. 2

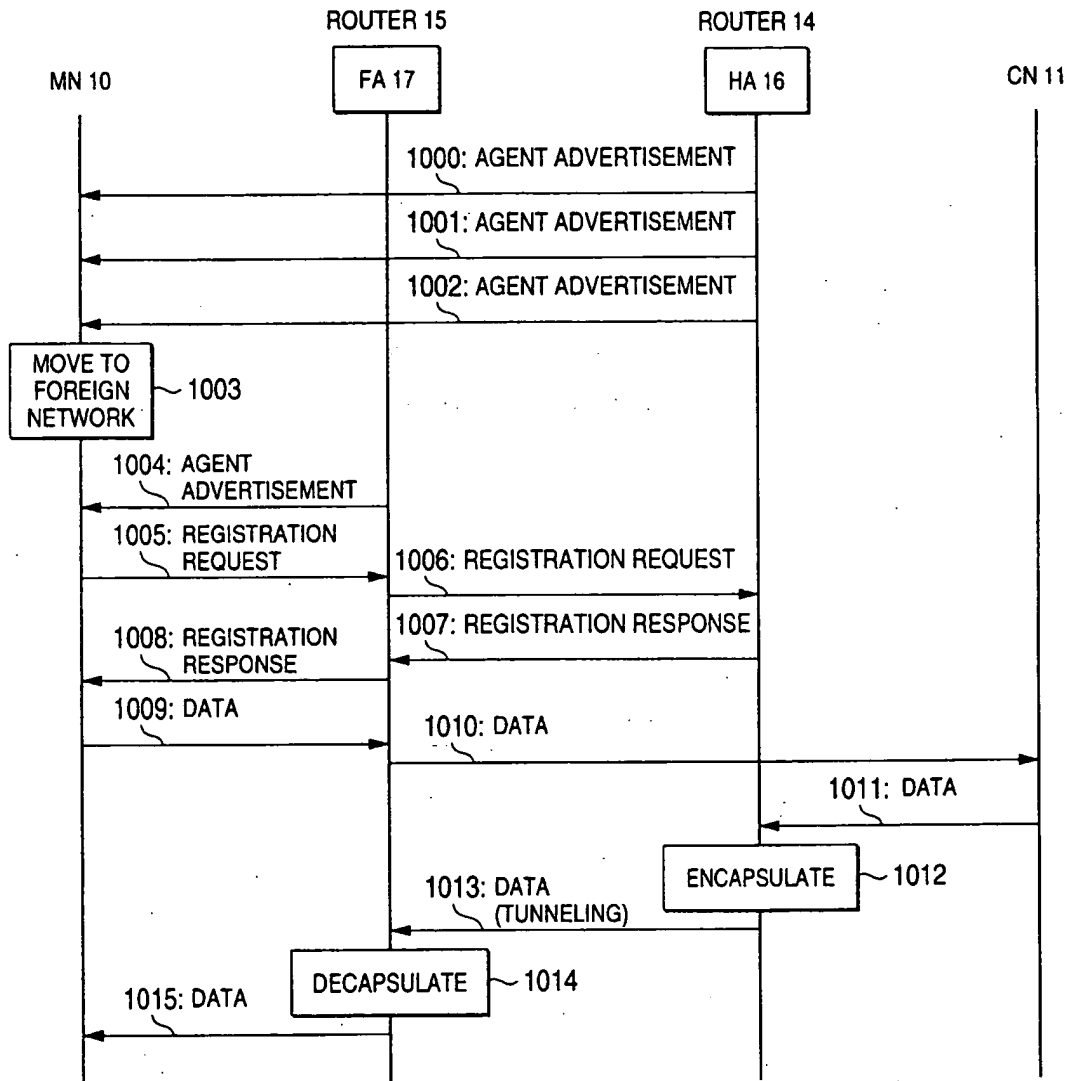


FIG. 3

TYPE	LENGTH	SEQUENCE NUMBER					
REGISTRATION LIFE TIME		R	B	H	F	M	G
							V
RESERVED							
ZERO OR CARE-OF ADDRESS							

FIG. 4

TYPE	S	B	D	M	G	Rsv	LIFE TIME
HOME ADDRESS							
HOME AGENT							
CARE-OF ADDRESS							
IDENTIFICATION							
EXTENSION ...							

FIG. 5

TYPE	CODE	LIFE TIME
HOME ADDRESS		
HOME AGENT		
IDENTIFICATION		
EXTENSION ...		

FIG. 6

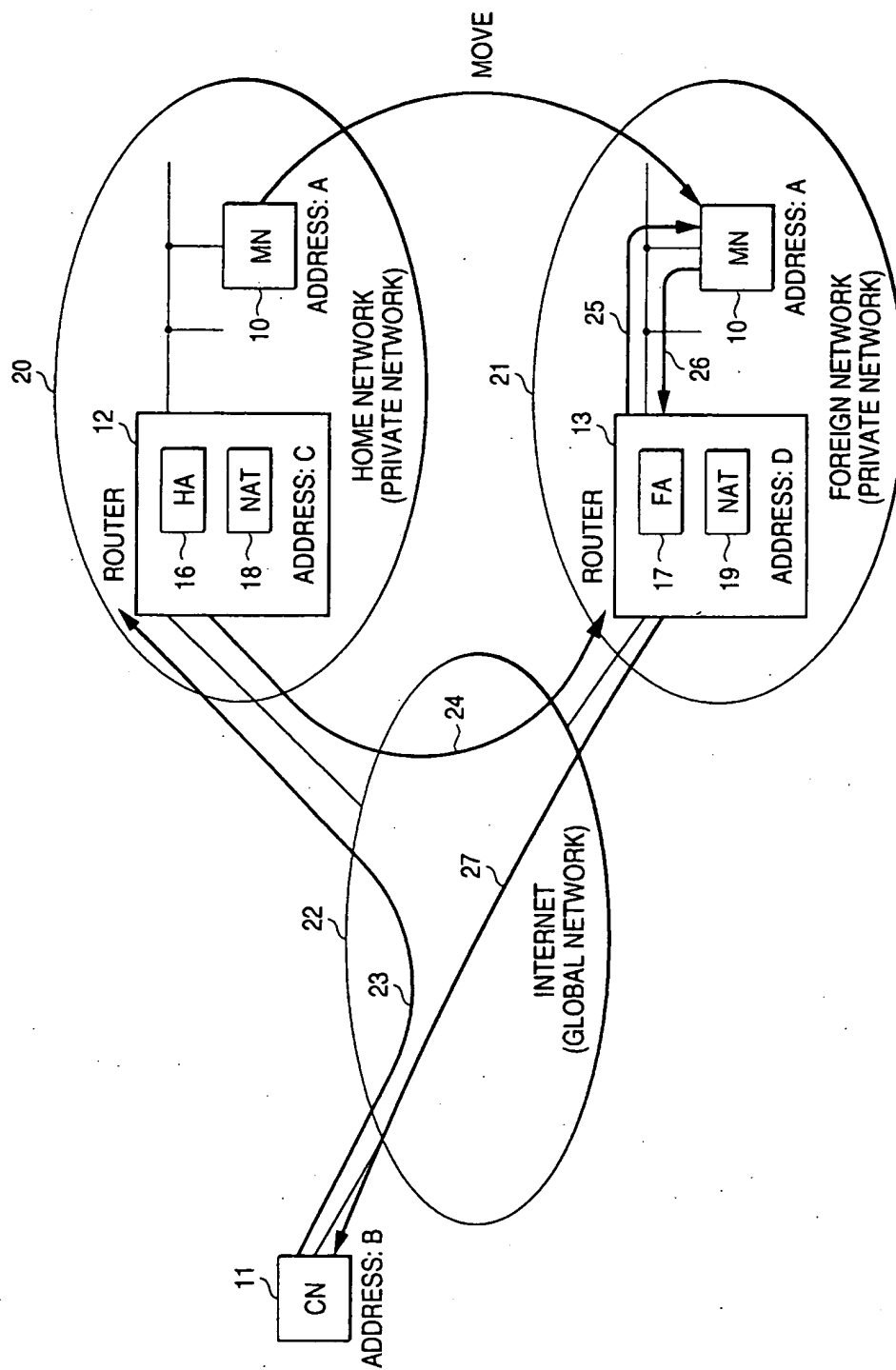


FIG. 7

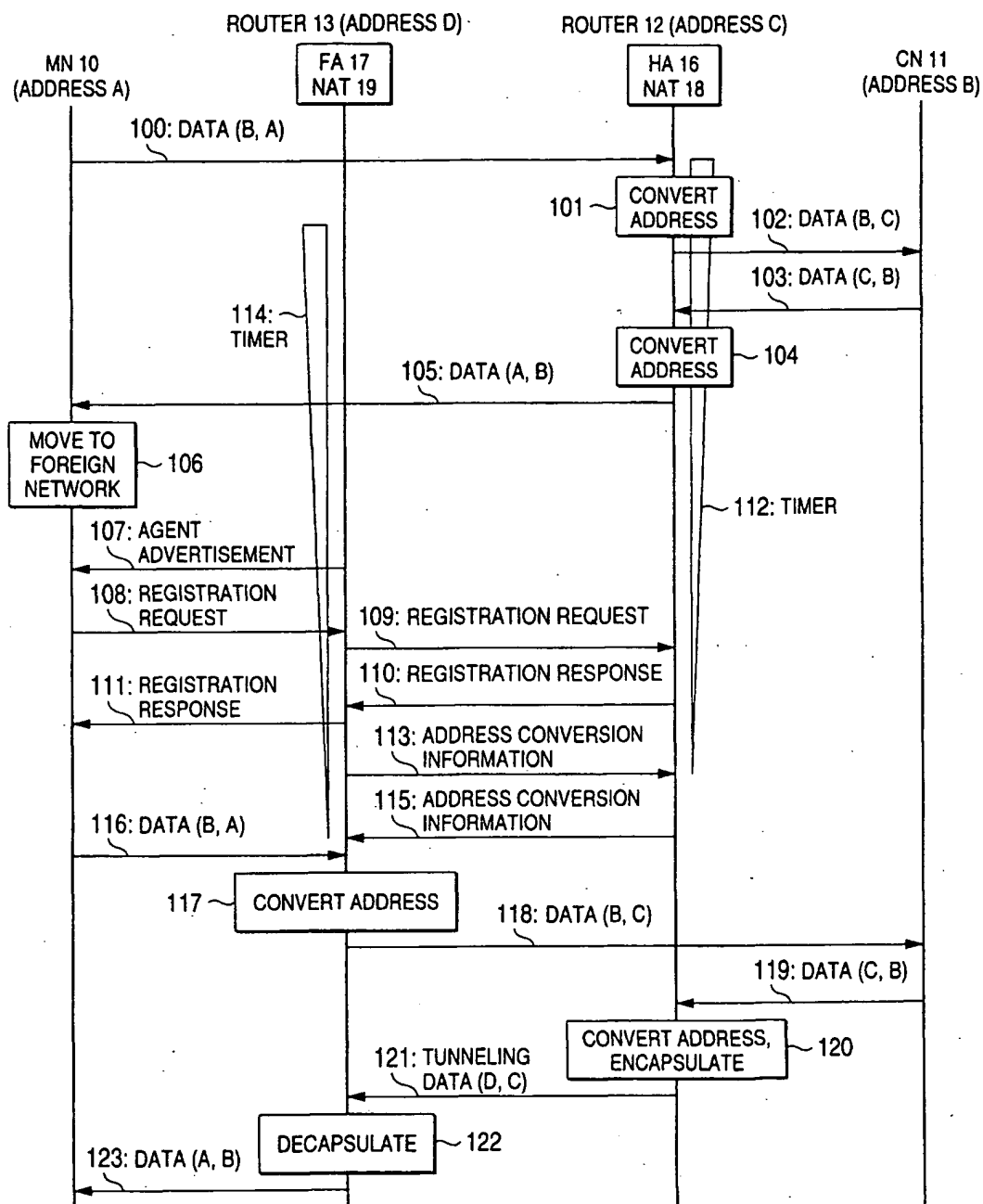


FIG. 8

28	HOME ADDRESS	CONVERTED ADDRESS	29
	A	C	
	⋮	⋮	

FIG. 9

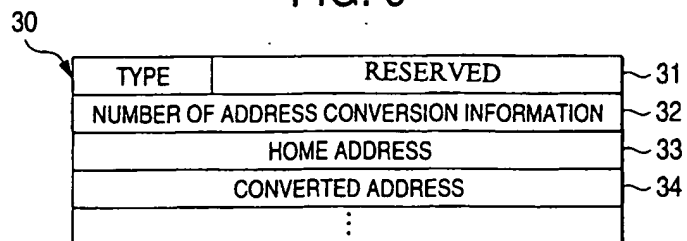


FIG. 10

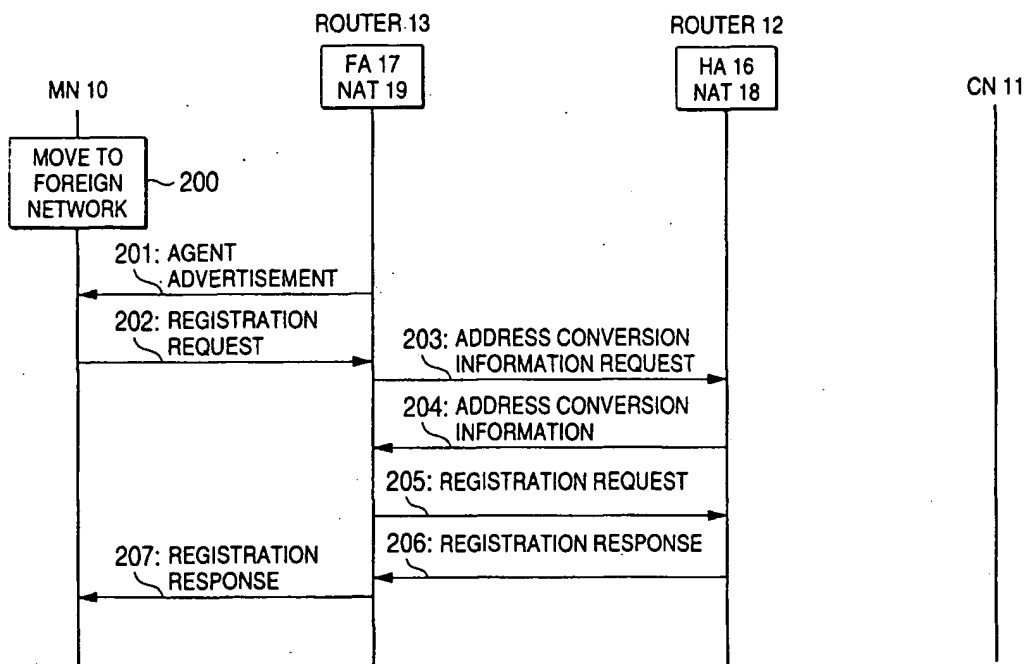


FIG. 11

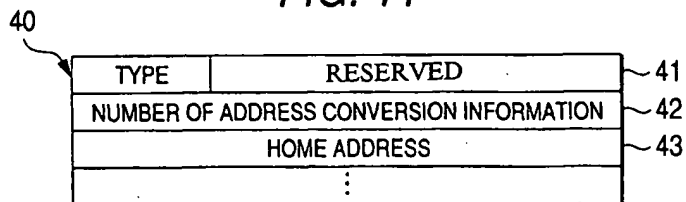


FIG. 12

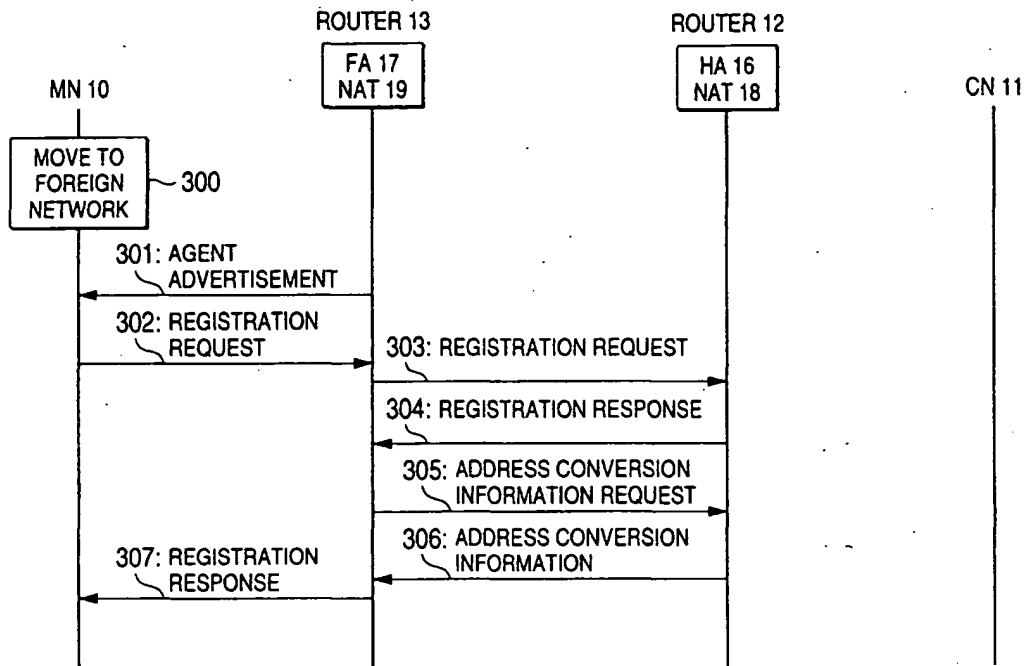


FIG. 13

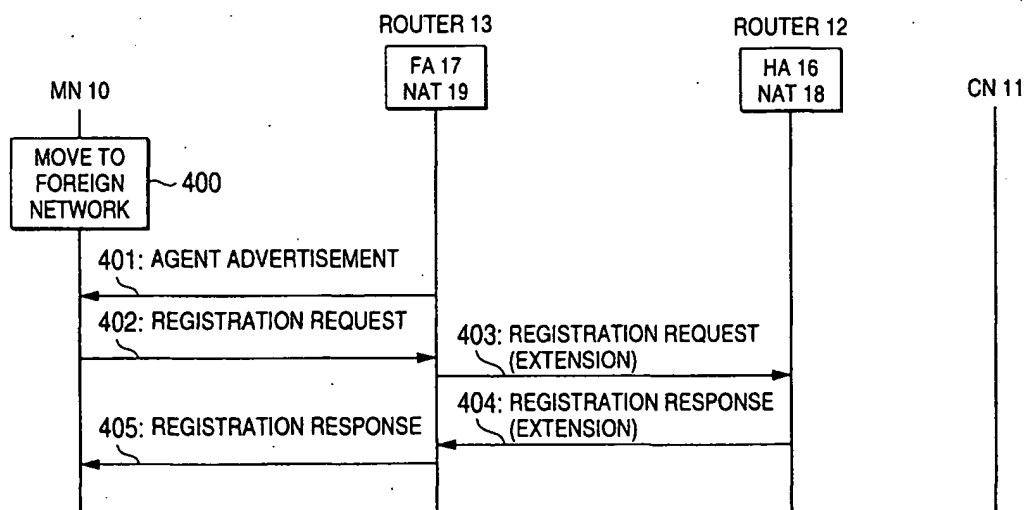


FIG. 14

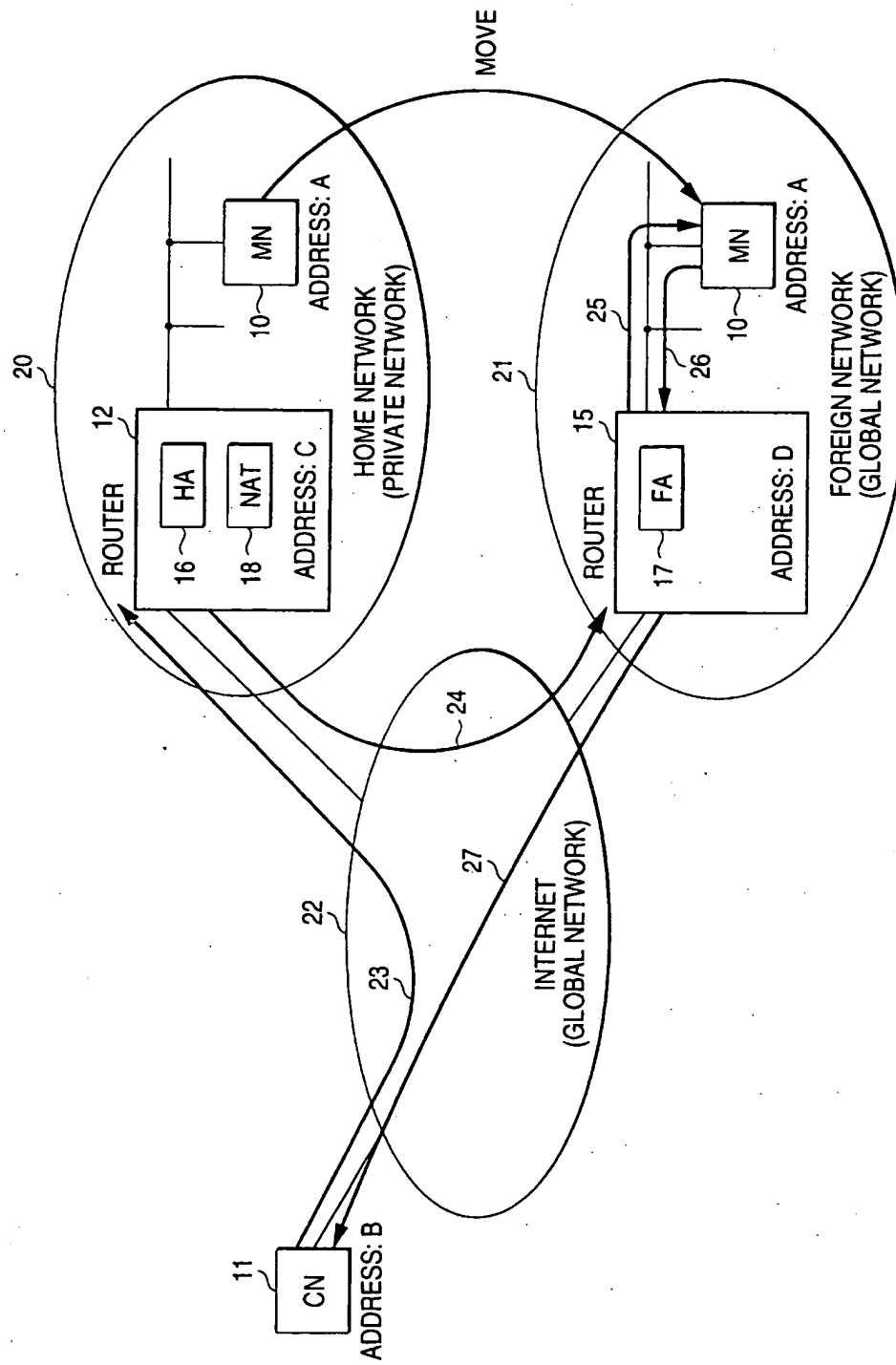


FIG. 15

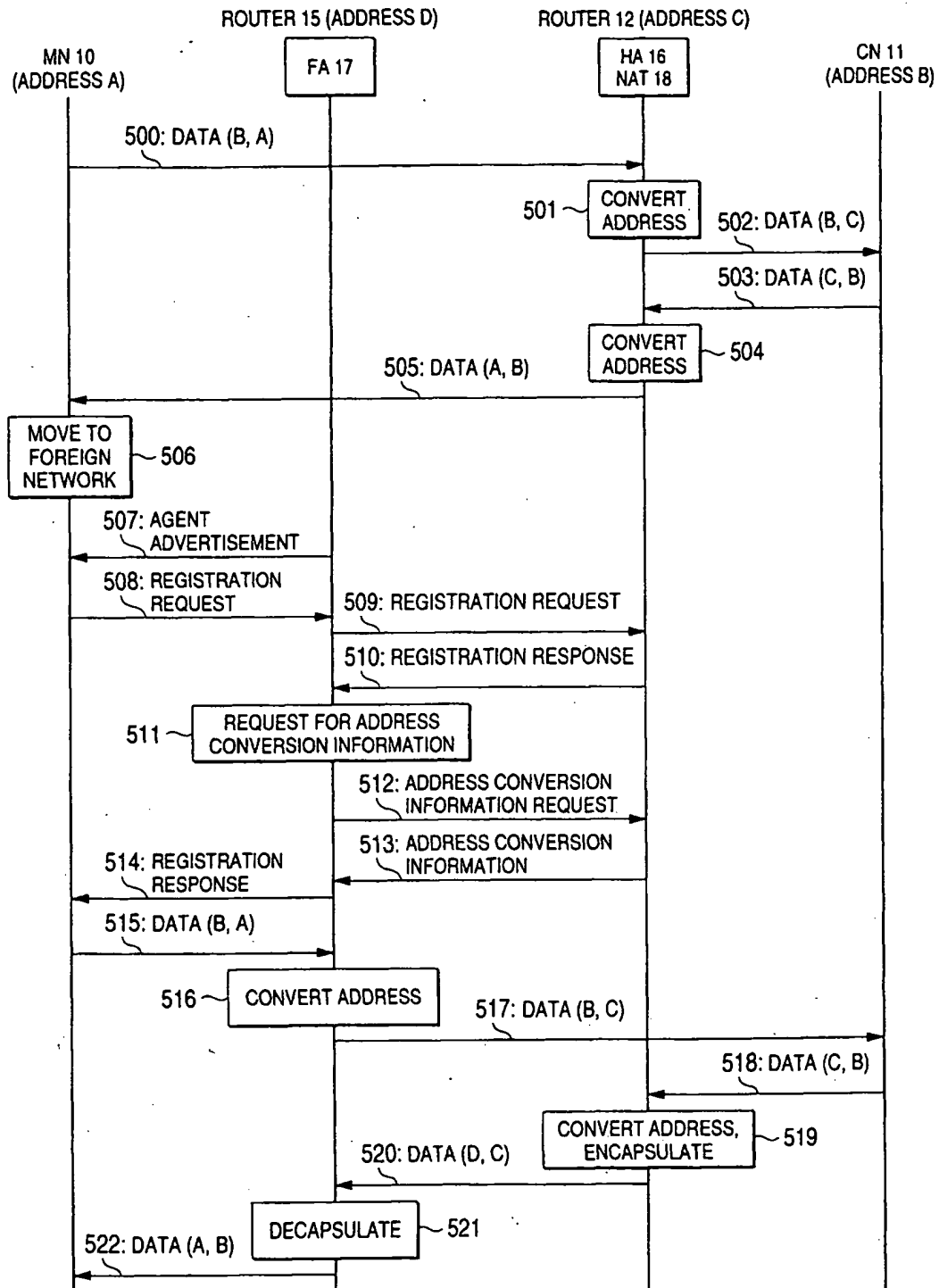


FIG. 16

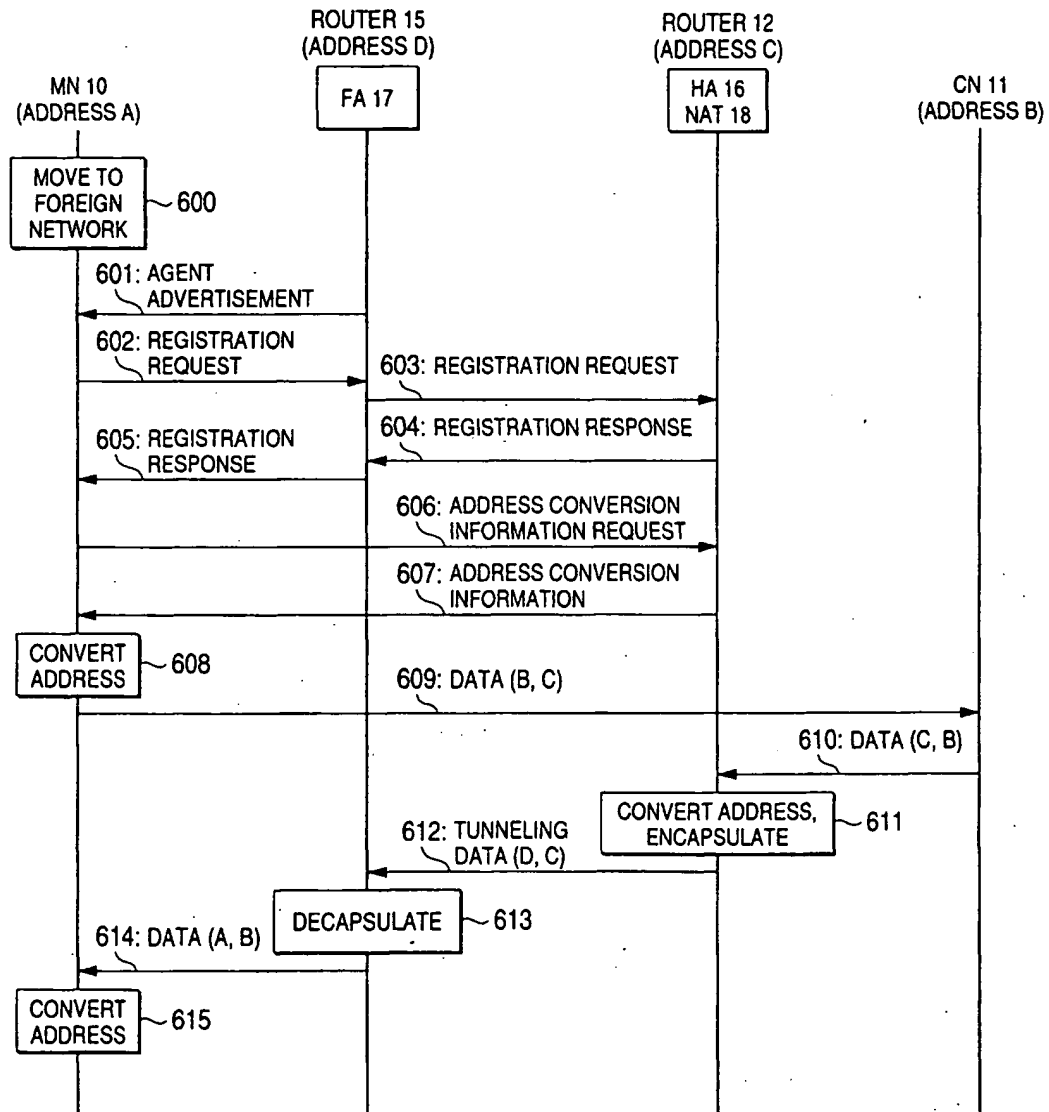


FIG. 17

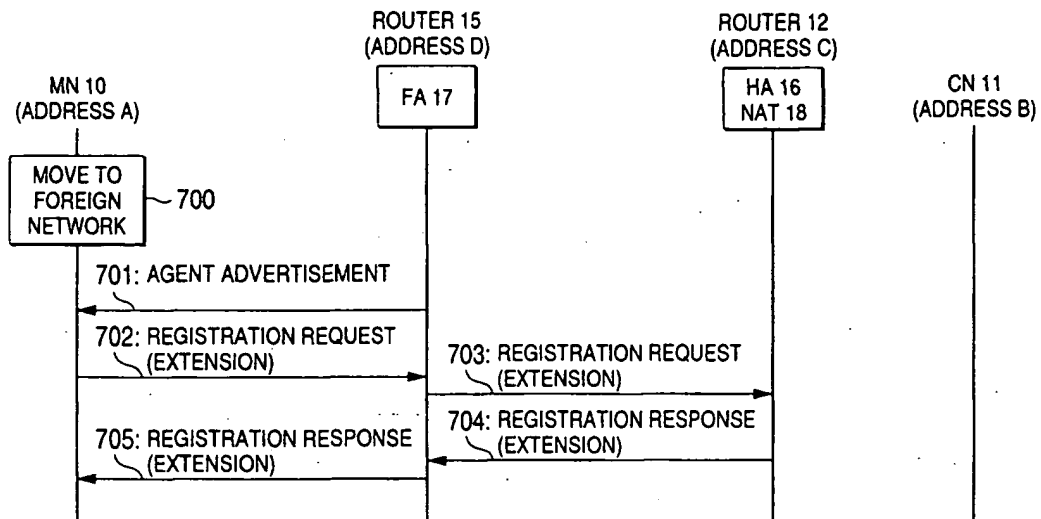


FIG. 18

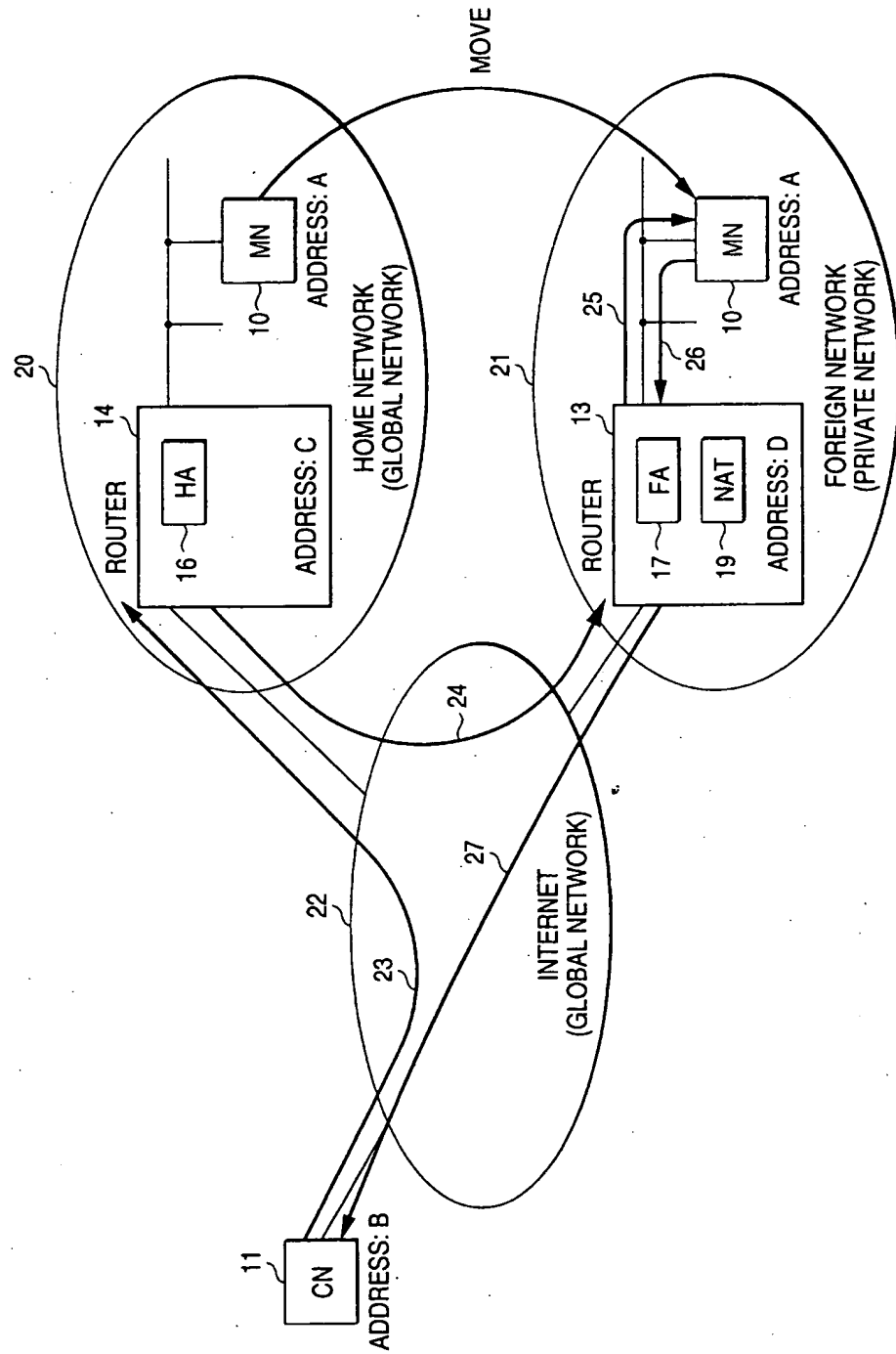


FIG. 19

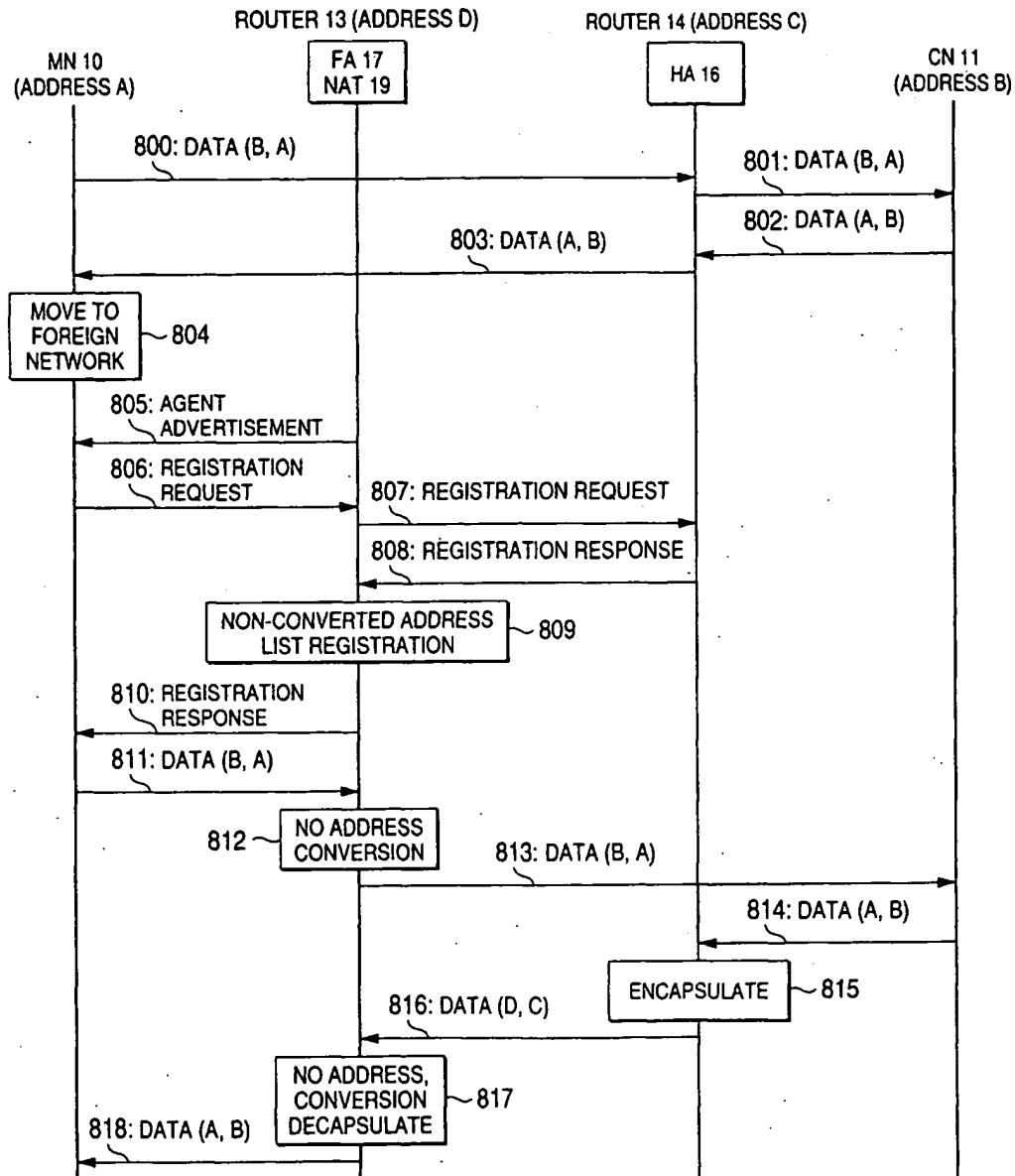


FIG. 20

NON-CONVERTED ADDRESS LIST
A
⋮

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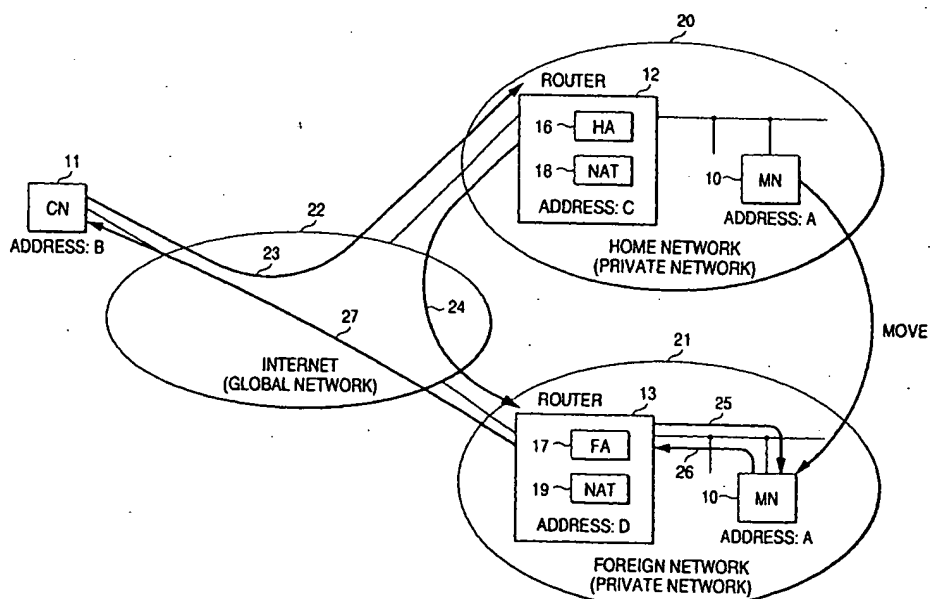
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(54) **Method for packet communication with mobile node**

(57) Even when a node (10) which can be moved among different networks (20, 21) and which sends a packet including a private address as a sender address sends the packet in either of the networks (20, 21), the sender address of the packet sent by the node (10) is

changed to a same global address. The method can prevent that a correspondent node (11) determines a packet from a same node as a packet from a different node due to a difference in their sender addresses, which hardly interrupts communication.

FIG. 6





European Patent
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EUROPEAN SEARCH REPORT

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EP 01 10 7526

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Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.7)
X	PATENT ABSTRACTS OF JAPAN vol. 1999, no. 02, 26 February 1999 (1999-02-26) & JP 10 303981 A (NEC CORP), 13 November 1998 (1998-11-13) * abstract *	1,2,7,9	H04L29/06
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			TECHNICAL FIELDS SEARCHED (Int.Cl.7)
			H04L
The present search report has been drawn up for all claims			
Place of search The Hague		Date of completion of the search 18 October 2004	Examiner Blanco Cardona, P
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**ANNEX TO THE EUROPEAN SEARCH REPORT
ON EUROPEAN PATENT APPLICATION NO.**

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This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report.
The members are as contained in the European Patent Office EDP file on
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18-10-2004

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For more details about this annex : see Official Journal of the European Patent Office, No. 12/82